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ParalLP: A Parallel Local Search Framework for Integer Linear Programming with Cooperative Evolution Mechanism

Peng Lin, Mengchuan Zou, Zhihan Chen, and Shaowei Cai*

Key Laboratory of System Software (Chinese Academy of Sciences) and State Key Laboratory of Computer Science,

Institute of Software, Chinese Academy of Sciences, Beijing, China

School of Computer Science and Technology, University of Chinese Academy of Sciences, Beijing, China

{linpeng, zoumc, chenzh, caisw}@ios.ac.cn

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Parallel Local Search for Integer Linear Programming

- Background and Motivation
- Parallel Framework for Solving ILP
- Experiments and Conclusions

Background

Let $m, n \in \mathbb{N}^+, A \in \mathbb{R}^{m \times n}, b \in \mathbb{R}^m, c \in \mathbb{R}^n$, and $l, u \in \mathbb{R}^n$. The optimization problem described by:

min
$$c^{\top} x$$

subject to: $Ax \leq b$
 $l \leq x \leq u$
 $x \in \mathbb{Z}^n$ (1)

is an instance of general integer linear programming (ILP).

- A fundamental model in operations research
- NP-Hard
- Strong descriptive capability









Background

Complete and Incomplete Algorithms

Complete Algorithms

- compute the optimal solution
- prove infeasibility

Branch-and-Bound



Academic Solver

- SCIP [Achterberg, 2009]
- HiGHS [Huangfu and Hall, 2018]

Commercial Solver

- Gurobi
- CPLEX

Incomplete Algorithms

• find high-quality solutions quickly

Local Search



Local-ILP [Lin et al., 2023]

Source code

https://github.com/shaowei-cai-group/Local-ILP

4 <u>P. Lin, et al</u>. Incomplete algorithms are important in industry applications of ILP

• The increasing computing power of multicore computer structures

 Although good serial local search algorithm has been proposed for ILP, the parallel local search has not been investigated

• Thus, we try to propose the first parallel local search framework for ILP

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Parallel Framework for Solving ILP



Figure 1: Our parallel local search framework for solving general integer linear programming with cooperative evolution mechanism.

How to generate high-quality initial solutions for each local search thread?

Definition 1. The *influence factor* of variable x_j in the objective function $\mathbf{c}^{\top} \mathbf{x}$, denoted as IO_j , is defined as

$$IO_{j} = \frac{(u_{j} - l_{j}) \cdot c_{j}}{\sum_{k=1}^{n} (u_{k} - l_{k}) \cdot |c_{k}|}$$
(2)

Similarly, the influence factor of variable x_j in constraint $A_i x \leq b_i$, denoted as IC_{ij} , is defined as

$$IC_{ij} = \frac{(u_j - l_j) \cdot A_{ij}}{\sum_{k=1}^n (u_k - l_k) \cdot |A_{ik}|}$$
(3)

Definition 2. Let m_j denote the number of constraints containing x_j , the **polarity** of variable x_j , denoted as P_j , is defined as

$$P_j = \frac{\sum_{i=0}^m IC_{ij}}{m_j} + IO_j \tag{4}$$

Polarity Initialization Polarity of variables (for high-quality) Random perturbations (for diversified solutions)

Manage a population balancing quality and diversity

The Fitness Function

 $R(\boldsymbol{s}) = R_Q(\boldsymbol{s}) \cdot p + R_D(\boldsymbol{s}) \cdot (1-p)$

The Objective Value Q(s) is initially set to -obj(s)

The Informative Degree:

- Whenever a solution's information is utilized, Q(s) is penalized
- If Q(s) > 0, Q(s) is updated as $Q(s) \times (1 \beta)$
- if Q(s) < 0, Q(s) is updated as $Q(s) \times (1 + \beta)$

The Differences from Other Solutions

$$D(\boldsymbol{s}) = \sum_{s' \in \text{Population}, s' \neq s} \sum_{j=1}^{n} |s_j - s'_j|$$

Generate new solutions from high-quality solutions in the population



- Background and Motivation
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ParalLP is significantly better than the state-of-the-art academic parallel solvers FiberSCIP and HiGHS

					:	#Fea	S				#Win										P(T)								
Benchmark Domain	#Ins	ŀ	HiGH	S	Fil	berS	CIP	ParaILP			HiGHS			FiberSCIP			ParaILP			HiGHS			FiberSCIP			ParaILP			
		10s	60s	300s	10s	60s	300s	10s	60s	300s	10s	60s	300s	10s	60s	300s	10s	60s	300s	10s	60s	300s	10s	60s	300s	10s	60s	300s	
Singleton	2	2	2	2	2	2	2	2	2	2	1	0	0	0	0	1	1	2	1	0.788	0.473	0.324	0.485	0.371	0.188	0.109	0.088	0.081	
Aggregations	2	0	1	1	1	1	1	1	1	1	0	0	1	1	1	0	0	0	0	1.000	0.615	0.523	0.672	0.531	0.508	0.761	0.633	0.555	
Bin Packing	2	0	2	2	0	1	1	1	2	2	0	2	0	0	0	1	1	0	1	1.000	0.905	0.715	1.000	1.000	0.993	0.979	0.923	0.773	
Equation Knapsack	3	0	0	0	0	0	0	0	0	1	0	0	0	0	0	0	0	0	1	1.000	1.000	1.000	1.000	1.000	1.000	1.000	1.000	1.000	
Knapsack	4	3	3	3	3	3	3	4	4	4	1	1	2	1	1	0	2	2	2	0.529	0.388	0.304	0.412	0.316	0.282	0.111	0.071	0.048	
Set Packing	5	2	4	4	3	4	4	4	4	4	1	1	1	0	0	1	3	3	2	1.000	0.802	0.609	0.759	0.549	0.323	0.369	0.260	0.219	
Cardinality	6	1	1	2	1	1	2	2	2	3	0	0	1	0	0	1	2	2	1	0.958	0.820	0.746	0.866	0.830	0.753	0.628	0.571	0.522	
Hybrid	7	3	3	4	2	4	4	4	5	5	0	0	1	0	0	0	4	5	5	1.000	1.000	1.000	1.000	1.000	0.699	0.689	0.548	0.524	
Mixed Binary	8	0	0	0	2	2	2	4	5	6	0	0	0	1	2	2	4	4	5	1.000	1.000	1.000	1.000	0.980	0.976	0.890	0.861	0.762	
Set Partitioning	9	1	2	4	3	4	6	4	6	7	0	0	0	0	2	3	4	4	4	1.000	0.999	0.927	0.930	0.884	0.805	0.887	0.827	0.765	
Set Covering	11	5	6	8	5	8	7	9	10	10	2	0	2	0	5	6	7	5	3	0.861	0.780	0.656	0.780	0.648	0.523	0.733	0.650	0.610	
Precedence	13	1	3	4	10	11	12	12	12	12	0	0	0	1	1	2	11	11	10	0.996	0.974	0.860	0.889	0.813	0.653	0.576	0.454	0.345	
General Linear	15	3	4	6	7	7	9	10	10	10	2	2	4	2	1	3	6	7	7	0.780	0.700	0.591	0.712	0.575	0.425	0.355	0.317	0.301	
Variable Bound	16	7	9	9	11	13	13	14	15	15	0	0	1	0	0	2	14	15	14	0.895	0.784	0.653	0.851	0.688	0.472	0.381	0.278	0.228	
Invariant Knapsack	18	4	7	11	10	11	13	12	13	13	2	1	3	2	3	2	12	12	11	0.924	0.847	0.736	0.778	0.706	0.664	0.549	0.475	0.421	
Total	121	32	47	60	60	72	79	83	91	95	9	7	16	8	16	24	71	72	67	0.911	0.831	0.730	0.818	0.724	0.610	0.582	0.508	0.452	

Table 1: Performance evaluation between SOTA academic solvers HiGHS, FiberSCIP and ParaILP.

ParalLP is competitive with the state-of-the-art commercial parallel solver Gurobi.

		#Feas													#Wir	ı				P(T)									
Benchmark Domain	#Ins	$Gurobi_{comp}$			$\operatorname{Gurobi}_{\operatorname{heur}}$			ParaILP			Gurobi _{comp}			$\operatorname{Gurobi}_{\operatorname{heur}}$			ParaILP			Gurobi _{comp}			$\operatorname{Gurobi}_{\operatorname{heur}}$			ParaILP			
		10s	60s	300s	10s	60s	300s	10s	60s	300s	10s	60s	300s	10s	60s	300s	10s	60s	300s	10s	60s	300s	10s	60s	300s	10s	60s	300s	
Singleton	2	2	2	2	2	2	2	2	2	2	0	2	1	1	0	1	1	1	0	0.237	0.113	0.042	0.240	0.115	0.045	0.109	0.088	0.081	
Aggregations	2	1	1	1	1	1	1	1	1	1	1	1	1	0	1	1	0	0	0	0.514	0.502	0.500	0.519	0.503	0.501	0.761	0.633	0.555	
Bin Packing	2	2	2	2	2	2	2	1	2	2	2	2	0	2	2	1	0	0	1	0.963	0.748	0.303	0.963	0.757	0.305	0.979	0.923	0.773	
Equation Knapsack	3	0	0	1	0	0	0	0	0	1	0	0	1	0	0	0	0	0	1	1.000	1.000	0.952	1.000	1.000	1.000	1.000	1.000	1.000	
Knapsack	4	3	3	4	3	3	4	4	4	4	1	2	2	1	1	4	2	2	1	0.317	0.271	0.100	0.315	0.274	0.093	0.111	0.071	0.048	
Set Packing	5	4	4	4	4	4	4	4	4	4	2	1	2	3	2	2	1	2	2	0.408	0.271	0.231	0.407	0.271	0.230	0.369	0.260	0.219	
Cardinality	6	2	3	3	2	3	3	2	2	3	1	2	1	0	1	3	1	1	0	0.741	0.543	0.364	0.741	0.543	0.359	0.628	0.571	0.522	
Hybrid	7	4	5	5	4	5	5	4	5	5	1	1	2	1	1	3	3	5	2	0.777	0.660	0.548	0.784	0.664	0.546	0.689	0.548	0.524	
Mixed Binary	8	2	3	3	2	3	3	4	5	6	1	0	1	1	1	1	4	4	5	0.986	0.975	0.960	0.986	0.970	0.958	0.890	0.861	0.762	
Set Partitioning	9	7	7	7	7	7	7	4	6	7	4	4	3	5	4	4	2	2	2	0.815	0.748	0.625	0.815	0.738	0.617	0.887	0.827	0.765	
Set Covering	11	9	9	9	9	9	9	9	10	10	5	5	4	2	6	7	4	2	2	0.616	0.384	0.291	0.613	0.371	0.264	0.733	0.650	0.610	
Precedence	13	12	12	12	12	12	12	12	12	12	4	3	4	1	4	9	8	6	2	0.684	0.529	0.350	0.688	0.534	0.361	0.576	0.454	0.345	
General Linear	15	9	11	11	8	11	11	10	10	10	6	5	6	3	8	8	4	3	3	0.421	0.286	0.251	0.460	0.295	0.254	0.355	0.317	0.301	
Variable Bound	16	14	14	14	13	14	14	14	15	15	2	4	7	2	5	5	12	10	8	0.612	0.430	0.297	0.595	0.399	0.291	0.381	0.278	0.228	
Invariant Knapsack	18	12	12	15	12	12	15	12	13	13	6	6	7	4	6	7	9	10	9	0.699	0.602	0.367	0.700	0.604	0.367	0.549	0.475	0.421	
Total	121	83	88	93	81	88	92	83	91	95	36	38	42	26	42	56	51	48	38	0.653	0.526	0.398	0.656	0.522	0.396	0.582	0.508	0.452	

Table 2: Performance evaluation between SOTA commercial solver Gurobi (both the exact and heuristic version) and ParaILP.

- The first parallel local search framework and an efficient parallel solver for solving general ILP.
- Our code: <u>https://github.com/shaowei-cai-group/ParaILP</u>
- Our framework could be easily extended by plugging other sequential local search algorithms as a subroutine to create new solvers.

Thank You! Q&A